

More is Merrier: Relax the Non-Collusion Assumption in Multi-Server PIR

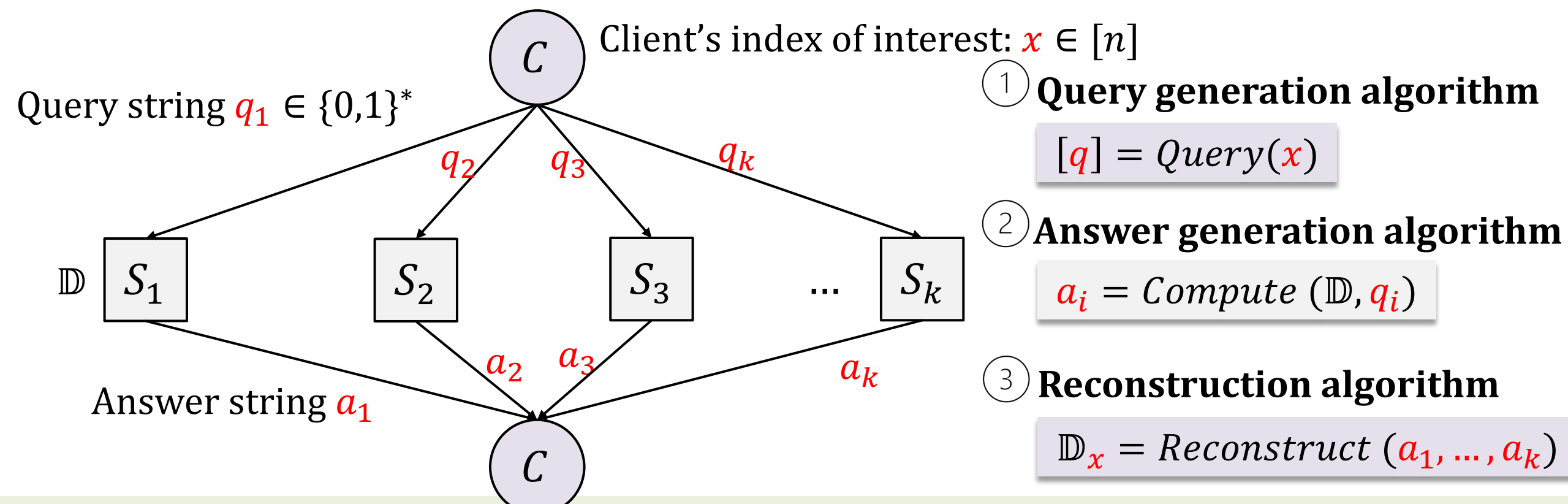
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What is Private Information Retrieval (PIR) [CKGS95]

A (k, t) -PIR scheme for a database $\mathbb{D}$ with n entries:

$\triangleright k$ -out-of- k , t -private
 $0 \leq t < k$



Security

- **Correctness** – C can reconstruct $\mathbb{D}_x$: $H(X|a_1 = \text{Compute}(\mathbb{D}, q_1), \dots, a_k = \text{Compute}(\mathbb{D}, q_k)) = 0$
- **Privacy** (IT, computational) – less than $t + 1$ parties learn no extra info: $H(X|\{q_j\}_{j \in S, |S| \leq t}) = H(X)$

$H(\cdot)$ computes the entropy of a random variable; X is the random variable for x

1. Multi-Server 1-Private PIR is Efficient

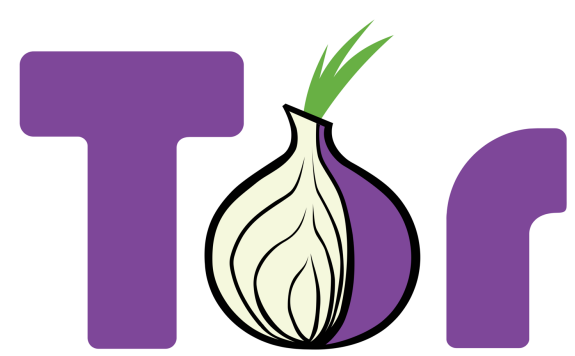
Our model: $(\ell, k, 1)$ -PIR where $\ell \geq k \geq 2$, database size $N = n \times b$
 $\triangleright \ell$ total servers, k -out-of- k , 1-private

Communication complexity	IT privacy	Computational privacy
Single-server	N	polylogarithmic
2-Server	$n^{O(\frac{\log \log n}{\log n})}$ [DG16]	polylogarithmic
3-Server	$2^{O(\sqrt{\log n \log \log n})}$ [Efremenko09]	-
k -Server	1-private: $n^{O(\frac{\log \log k}{k \log k})}$ [BIKRO2] t -private: $O(\frac{k^2}{t} \log(k n^{1/(2k-1)}))$ [WY05]	1-private: $(\lambda + 2) \log k \left\lceil \log \frac{n}{\lambda} \right\rceil + \frac{\ell}{\ell-1} b$ [HH19] (e.g., $\lambda = 128$, then $130 \log k (\lceil \log n \rceil - 7) + \ell/(\ell-1)b$)
$O(\log n)$ -Server	$O(\log^2 n \log \log n)$ [CKGS95]	-

Computation complexity (cPIR): polylogarithmic [LMW23] VS $\frac{N}{\ell-1}$ bit operations per server [HH19]

But $(\ell, k, 1)$ -PIR assumes **no pairwise collusion**

yet collusion is **easy** to implement over unobserved communication channels and **impossible** to detect



Anonymous channels



Encrypted communication channels



Side channels

Can we **relax the non-collusion assumption**?



2. What we do and How

What : Relax the non-collusion assumption to **rationality** assumption, i.e., servers are rational or malicious

How: Given two premises / our facts

1. We cannot directly detect collusion and collusion can be realized with any protocol, e.g., MPC, TEE, magical clouds, etc.
2. After successful collusion, at least some colluding parties have learned something nontrivial about the index x or entry $\mathbb{D}_x$ --- denote as $f(x)$

we **design a mechanism** such that

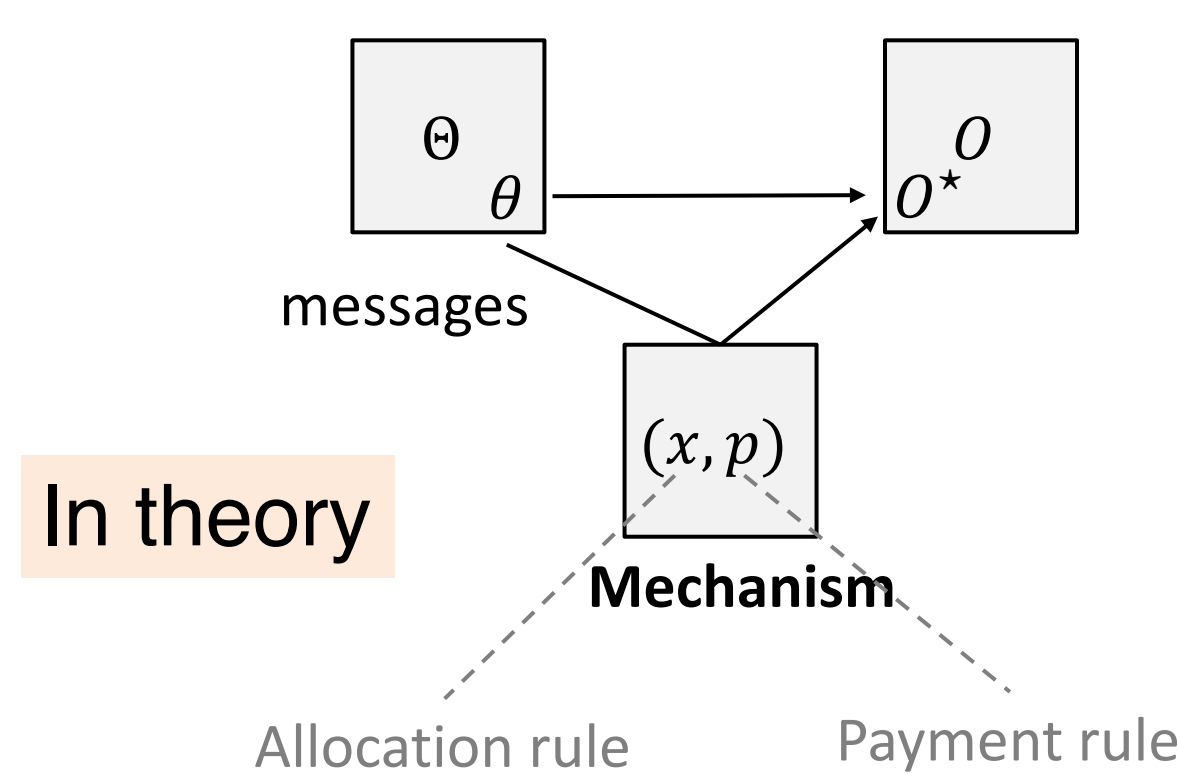
- (a) it induces a **game** where **exactly** one of the servers can take advantage of the **information gain** $f(x)$ to maximize its utility at the expense of others,
- (b) resulting in some party **unwilling** to collude to give others such an advantage

What is a mechanism

Agents' type space

Outcome space

An **algorithm** on a public bulletin board that participants can interact with



In theory

In practice



3. Mechanism overview

Mechanism M

Unknown: secrets $f(x), f(x^r)$

Known: server set $\{S_1, \dots, S_\ell\}$

$\triangleright$ Winner selection rule W

If any server S_i tells M the correct secret **first** along with its input and a **zero-knowledge proof** [Groth16, GWC19, OB22] of inequality is not provided by other servers S_{-i} by time Δ , select S_i as **winner** and mark all other servers as **colluders**

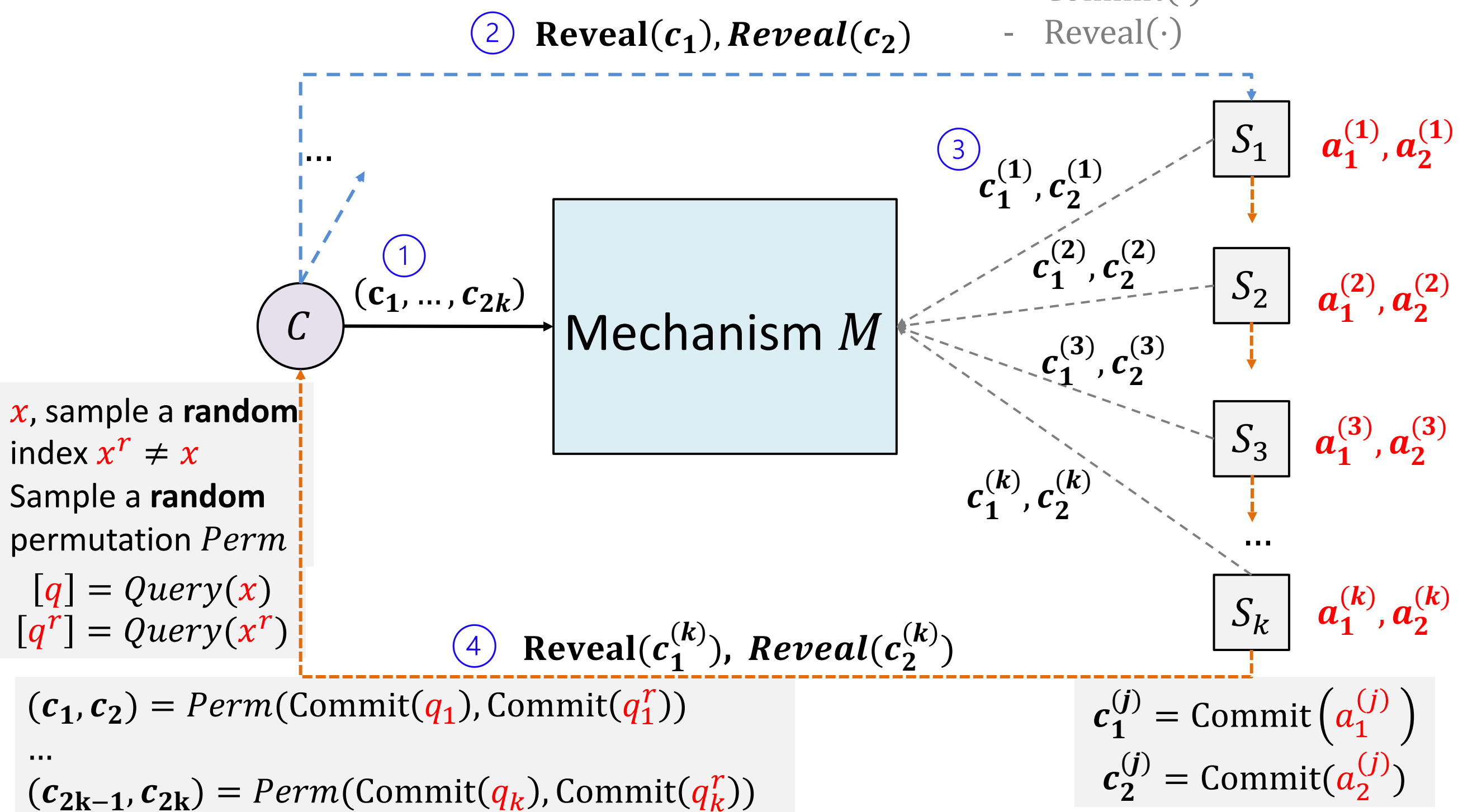
$\triangleright$ Payment rule P

- (1) Reward the **winner** a proper amount;
- (2) Penalize each marked **colluder** a proper amount;
- (3) Penalize S_1 a proper amount if it tells a wrong secret;
- (4) Charge proper amount of service fees for each queried server from the client and transfer to servers if there is no collusion after a privacy protection window

4. Protocol overview

Assume a secure commitment scheme:

- Commit($\cdot$)
- Reveal($\cdot$)



x , sample a **random** index $x^r \neq x$
Sample a **random** permutation Perm

$[q] = \text{Query}(x)$
 $[q^r] = \text{Query}(x^r)$



5. Communication and computation overhead

On paper One additional commitment per message – instantiated with SHA-3

Implementation as a smart contract on Ethereum

CheckCircuits($\cdot$) checks if the function is trivial with oracle services

Table 1. Gas costs summary

Normal service	Gas	Dollars	Collusion resolution	Gas	Dollars
Contract deployment	4697299	\$8.63	Accuse($\cdot$)	223766	\$0.41
Deposit($\cdot$)	105436	\$0.19	CheckCircuits($\cdot$)	66991+	\$0.12+
PostRequests($\cdot$)	405657	\$0.74	VerifyExchange($\cdot$)	61822	\$0.11
SubmitResponse($\cdot$)	97400	\$0.18	VerifyGeneralFunc($\cdot$)	275279	\$0.51
ClaimServiceFee($\cdot$)	33103	\$0.06	zkVerify($\cdot$)	2286423	\$4.20