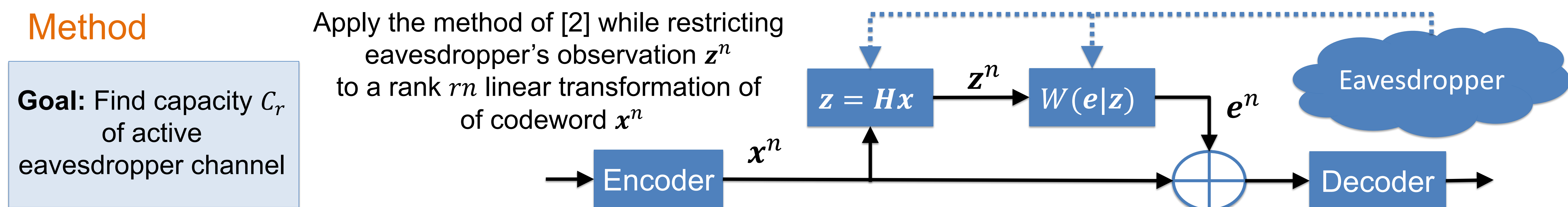
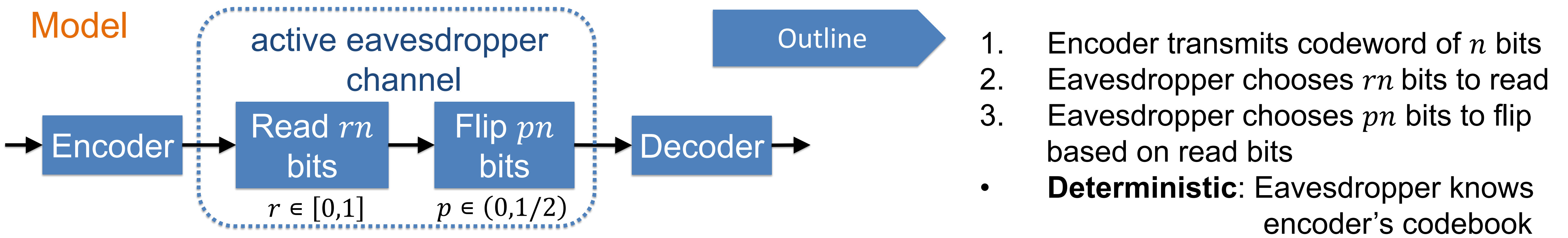


# A Characterization of the Deterministic Capacity of a Channel with an Active Eavesdropper

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**Overview** We investigate a family of channels in which an active eavesdropper may read and flip a fraction of transmitted bits. This active eavesdropper channel is sometimes referred to as a limited view (LV) adversarial channel [1]. We characterize the capacity of the channel family in a *deterministic* setting, i.e., when no private randomness is shared between the encoder and decoder.



## Results & Conclusions

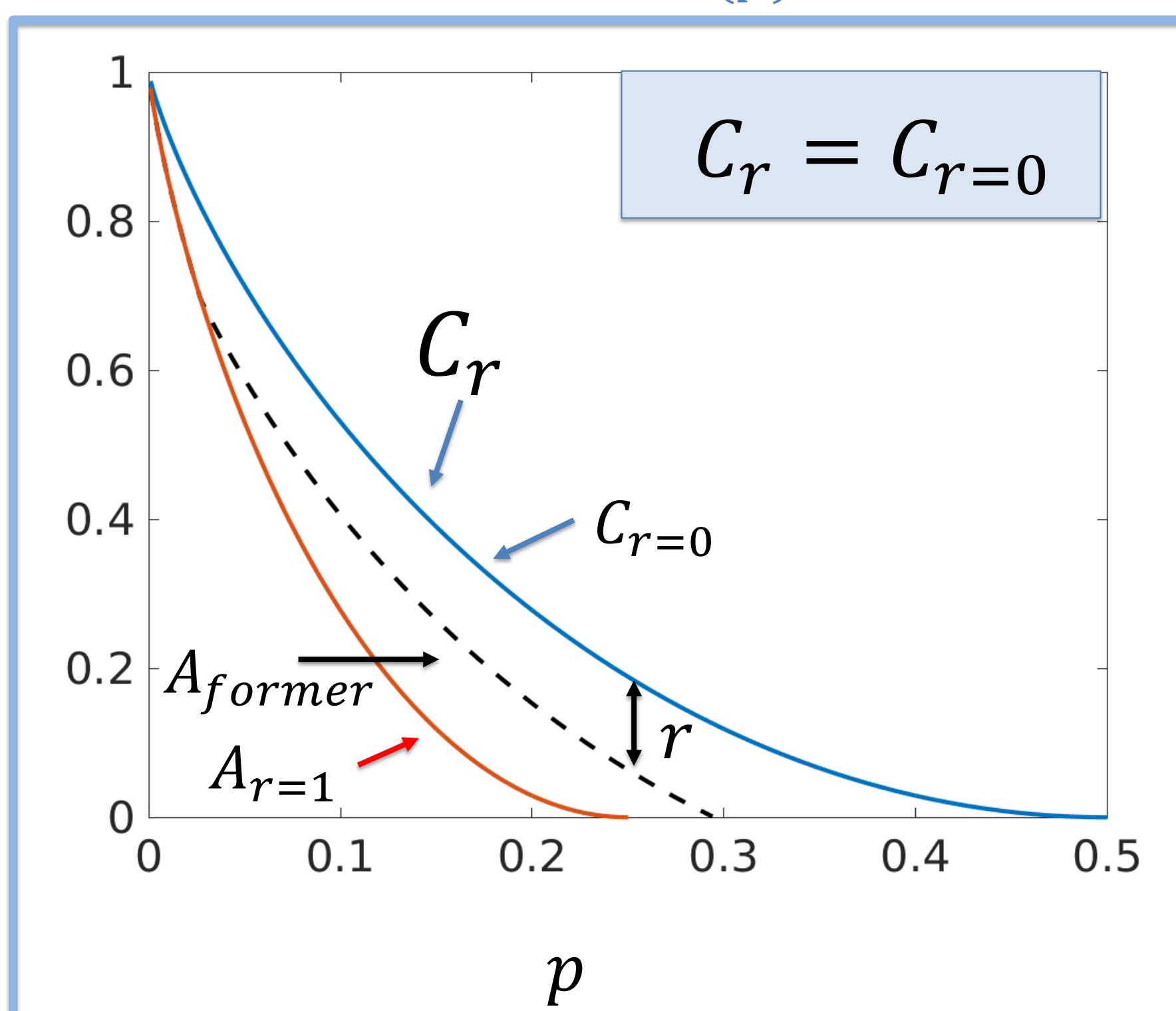
Capacity

Consider values of  $r \in [0, 1]$  partitioned into 3 regions

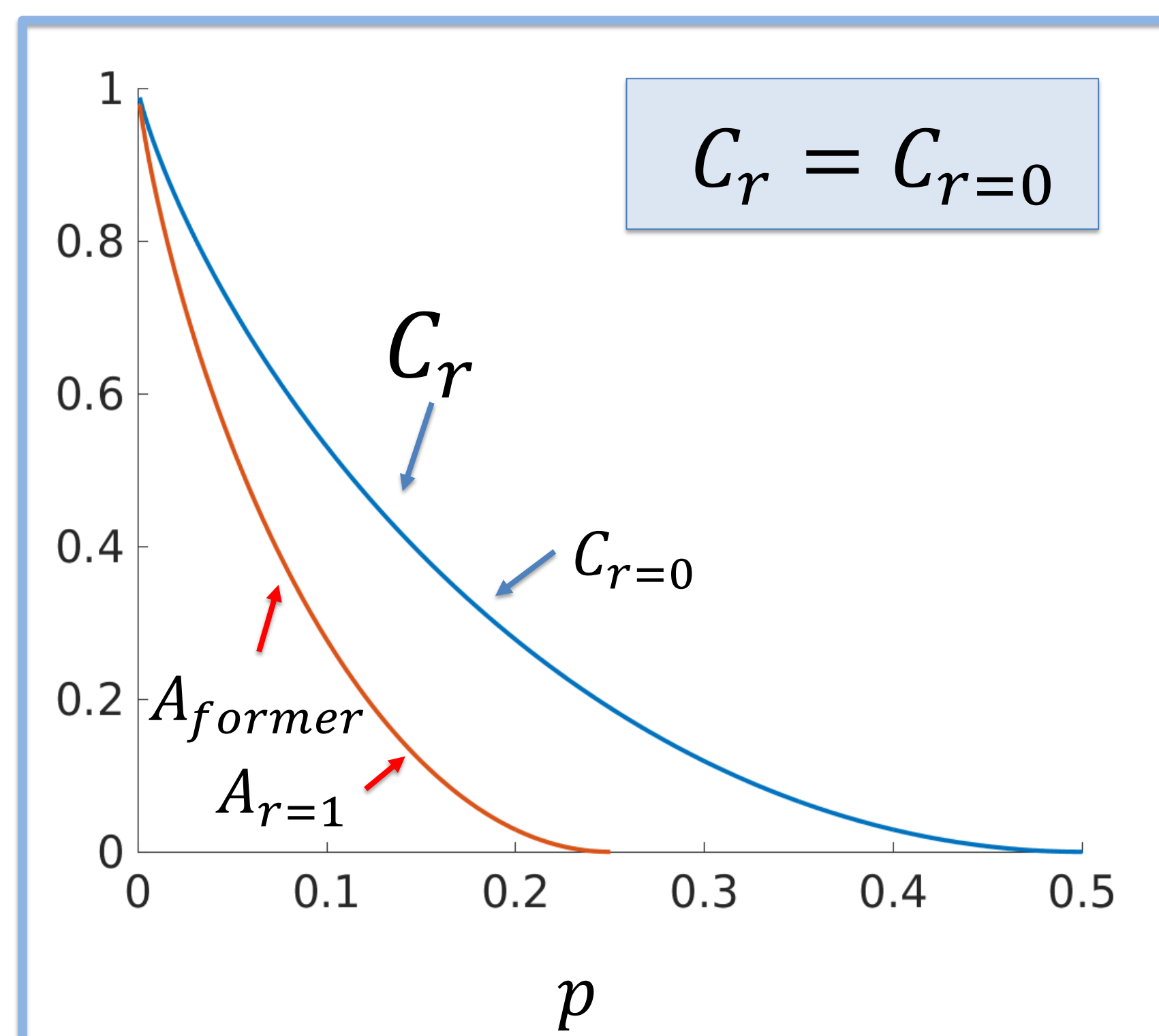
Key

- $C_r$  := capacity of active eavesdropper channel
- $A_r$  := lower bound on achievable rate of eavesdropper channel
- $A_{former}$  := former largest known lower bound on achievable rate [2]
- Capacity  $C_{BSC(p)}$  of binary symmetric channel (w/ parameter  $p$ )
- Gilbert-Varshamov Bound

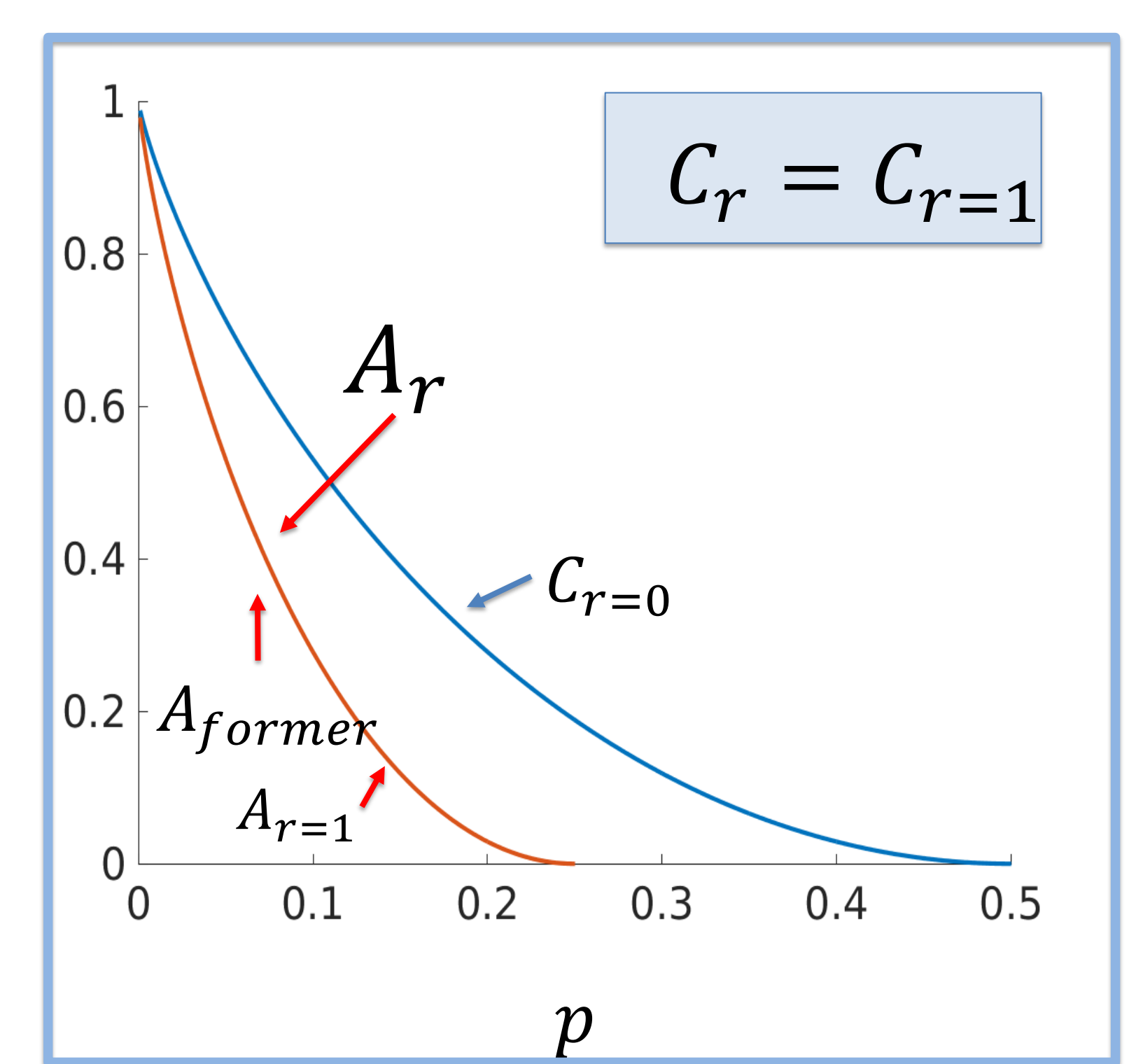
If  $r \in [0, C_{BSC(p)}/3]$



If  $r \in [C_{BSC(p)}/3, C_{BSC(p)}]$



If  $r \in [C_{BSC(p)}, 1]$



Eavesdropper's Optimal Strategy

- When  $r$  is large, decode codeword and push to nearest neighbor
- When  $r$  is small, pick flipped bits randomly (independent of codebook/codeword)
- Always pick readable bits randomly (independent of codebook)

Takeaway

**A little eavesdropper uncertainty (i.e.,  $r < 1$ ) can dramatically expand the achievable rate region**

## Ref

- [1] Wang and Safavi-Naini, "Limited view adversary codes: bounds, constructions and applications" ICITS 2015
- [2] Langberg, "Oblivious communication channels and their capacity", IEEE TOIT 2008