

## HexCFI: Context Sensitive Dynamic Control-Flow Integrity

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### Motivation

- Despite hardening with LLVM-CFI, Chrome is still vulnerable to control-flow hijacks
- Attacks target *indirect* control flow transfers that are computed at runtime
- LLVM-CFI *statically* computes allowed target sets for indirect control-flow transfers and is **over-approximate**

### Motivating Example

```
void foo() { }
void bar() { }
void fun() { }

int main( ) {
    void(*fnptr) ();
    int a = 2;
    if(a % 2 == 0)
        fnptr = &bar;
    else
        fnptr = &foo;
    fnptr();
    return 0;
}
```

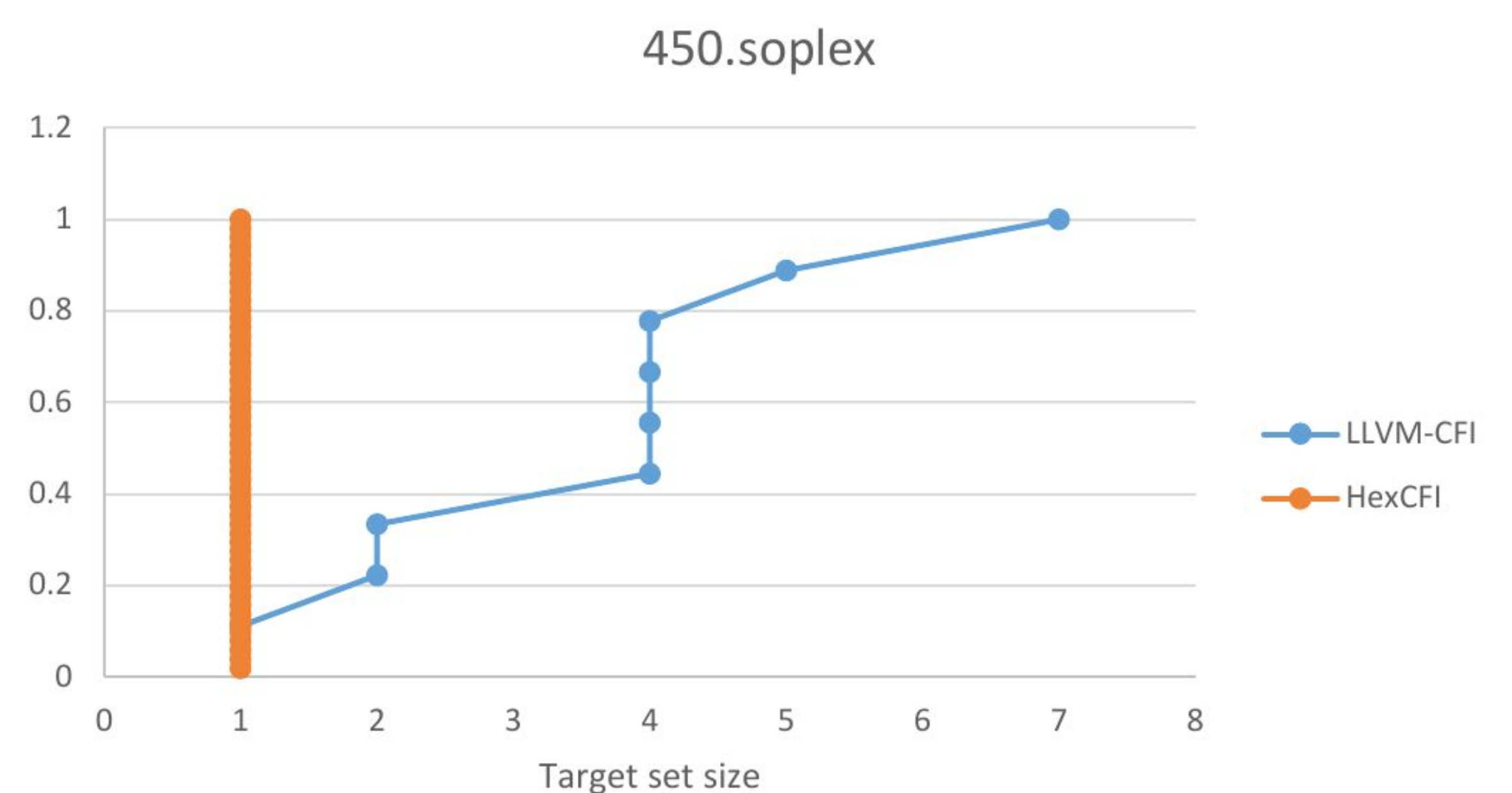
#### Target Sets

- LLVM\_CFI { foo(), bar(), fun() }
- HexCFI { bar() }

### Goals

- Compute an **optimum** target set per indirect call site
- An optimum target set is the smallest set such that the program can still execute correctly
- This will provide the strongest possible CFI security policy
- Promote call sites with one target to direct calls

### Evaluation



## HexCFI Architecture

#### Analysis Phase

- Instrumentation logs targets for each indirect control-flow transfer
- Test Suite / fuzzing used to observe all valid execution paths

#### Enforcement Phase

- Computes target set per callsite from Target Log
- Instruments indirect call sites to enforce valid target set

