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The Center for Education and Research in Information Assurance and Security

HexVASAN: A Sanitizer for Variadic Functions

hexhive

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Motivation

➢ Functions that take variable number of arguments are not checked statically, e.g.,
  `int printf(const char*format,…);`

➢ Vulnerabilities arise from differences between how the caller passes the variadic arguments and how the callee uses them.

➢ HexVASAN analyzes variadic functions and enforces runtime integrity checks.
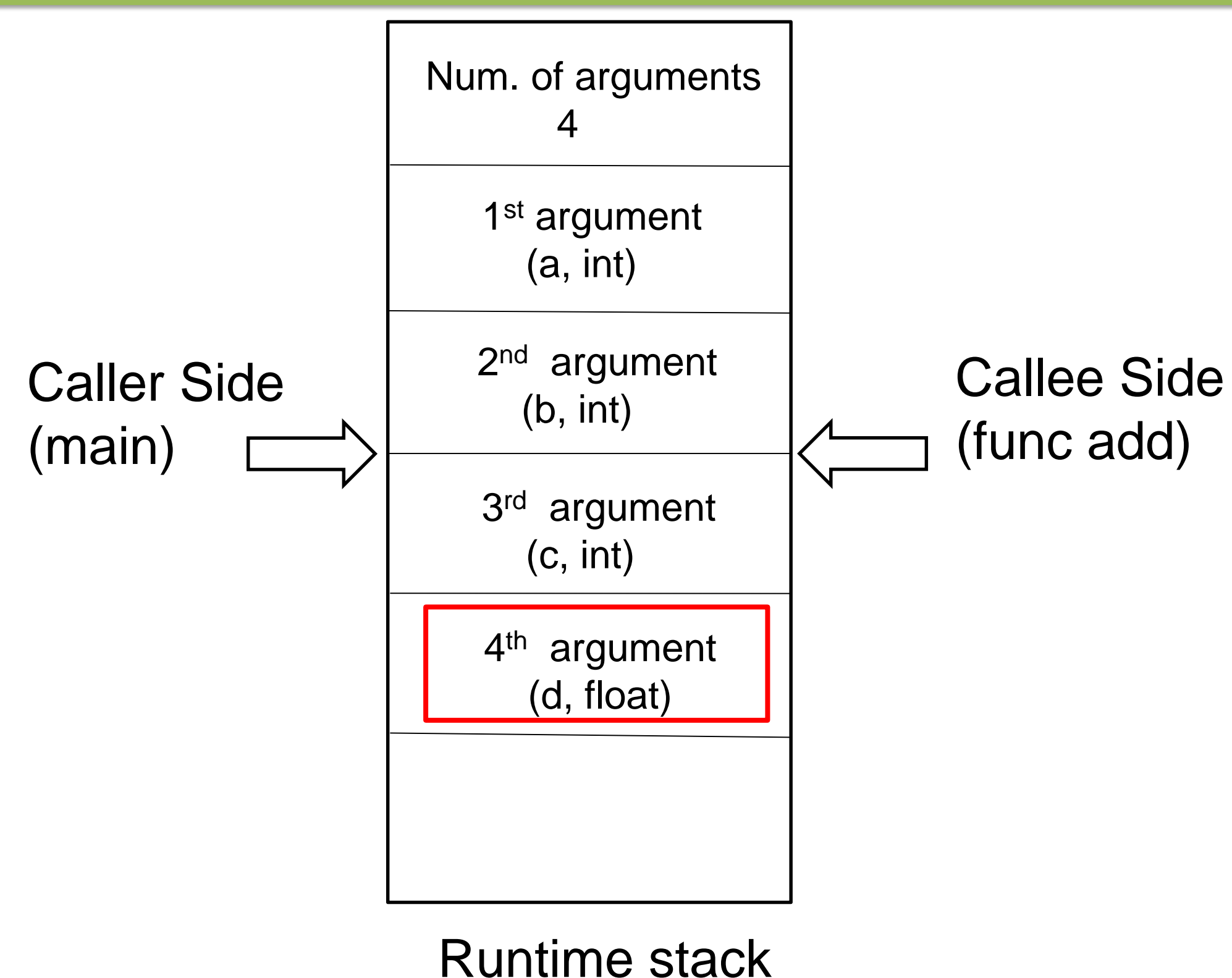
Attack Model

```
printf("%200$p%n", &var)
```

➢ In a format string attack, an attacker controls the first argument and changes how many parameters are used and interpreted.
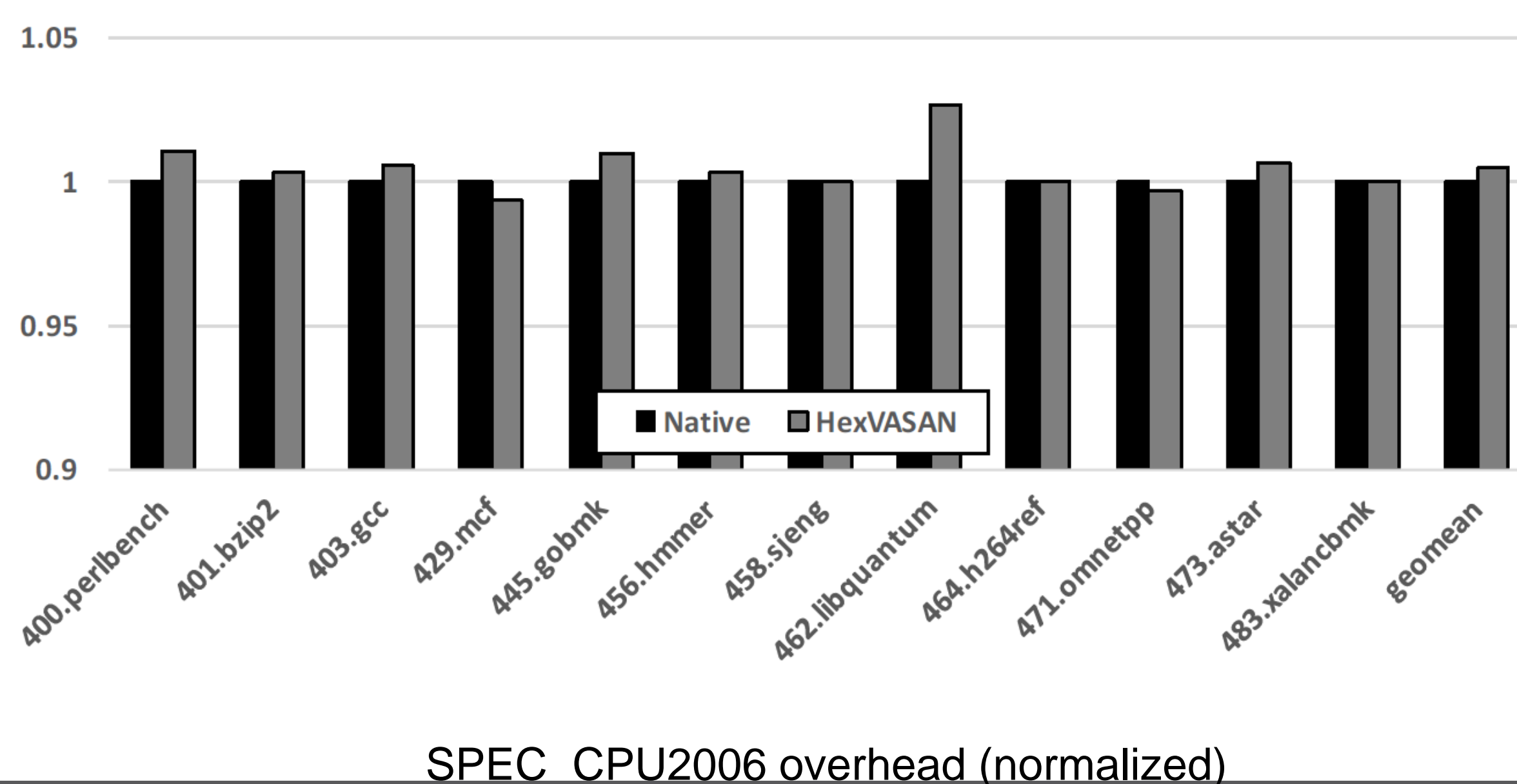
➢ HexVASAN prevents such attacks.

HexVASAN Architecture

```
int add(int n, ...) {
    va_list list;
    va_start(list, n);
    for (int i=0; i < n; i++) {
        total = total + va_arg(list, int);
    }
    va_end(list);
    return total;
}

int main(int argc, const char *argv[]) {
    int a, b, c, result;
    float d;
    result = add(4, a, b, c, d);
    return 0;
}
```

| Num. of arguments 4 |
| 1st argument (a, int) |
| 2nd argument (b, int) |
| 3rd argument (c, int) |
| 4th argument (d, float) |

Caller Side (main) →   ← Callee Side (func add)

Runtime stack

Result

SPEC CPU2006 overhead (normalized)

Conclusion

➢ Real software such as Firefox, Chromium uses variadic functions both directly and indirectly, leading to a potential attack surface.

➢ HexVASAN successfully tracks if there is any type mismatch and prevents exploits.

➢ Incurs 0.45% and 1% overhead for SPEC CPU2006 and Firefox, respectively.

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