

Privacy-Preserving and Efficient Friend Recommendation in Social Networks

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Introduction

Friend recommendation is a well-known application in many social networks and has been studied extensively in the recent past. However, with the growing concerns about users' privacy, there is a strong need to develop privacy preserving friend recommendation methods for social networks. In this paper, we propose two novel methods of **Privacy-Preserving Friend Recommendation (PPFR)** based on common neighbors proximity. The first method is based on the properties of **additive homomorphic encryption** scheme and also utilizes a universal hash function for efficiency purpose. Nevertheless, this efficiency comes at the expense of degraded accuracy due to the involved hash collisions. Whereas, the second method utilizes the concept of **protecting the source privacy** through randomizing the message passing path and recommends friends accurately.

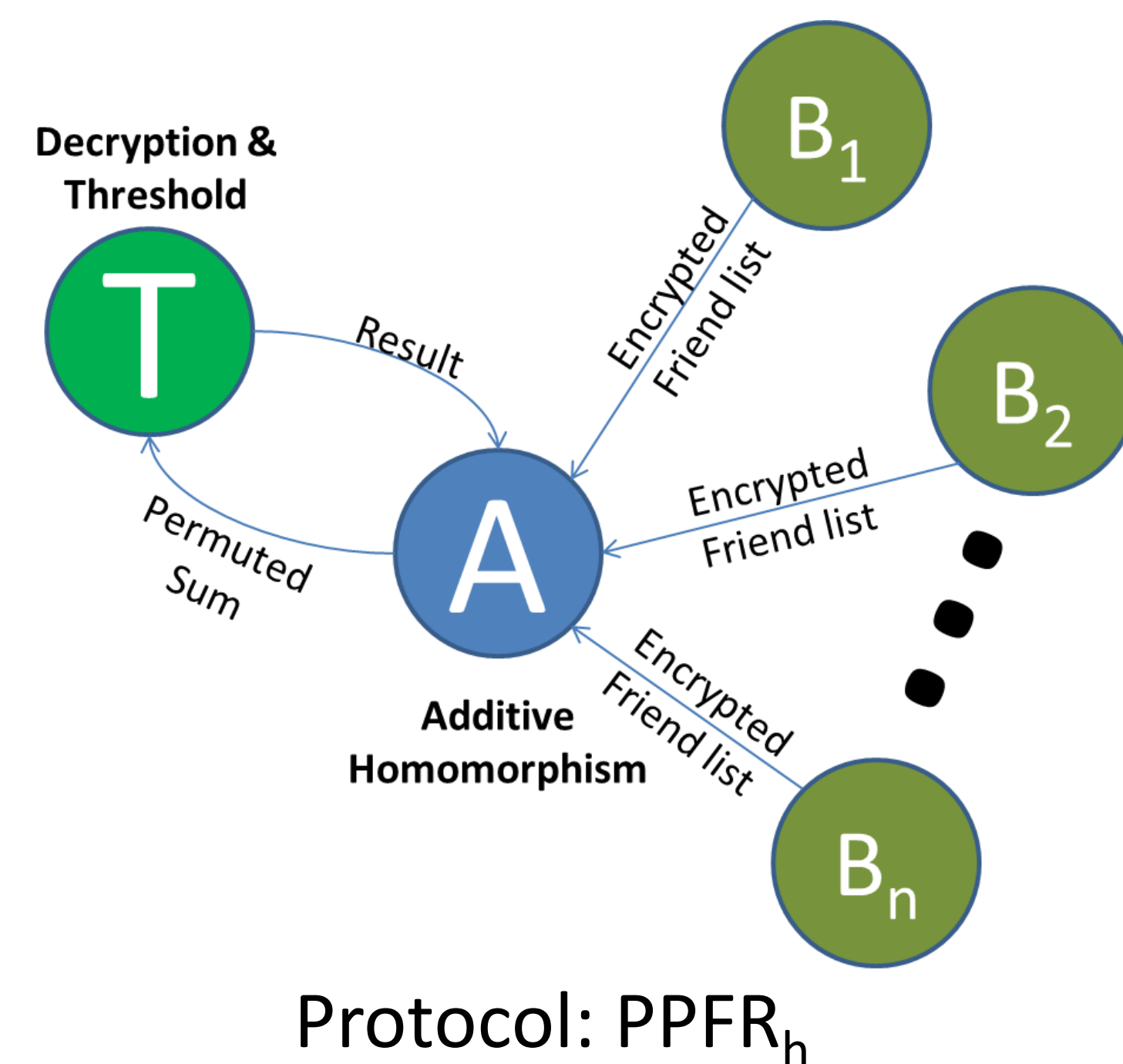
Two Protocols

The protocol $PPFR_h$ utilizes the, **additive homomorphic encryption scheme**[1] to sum up the friend list of all A's friends under encryption.

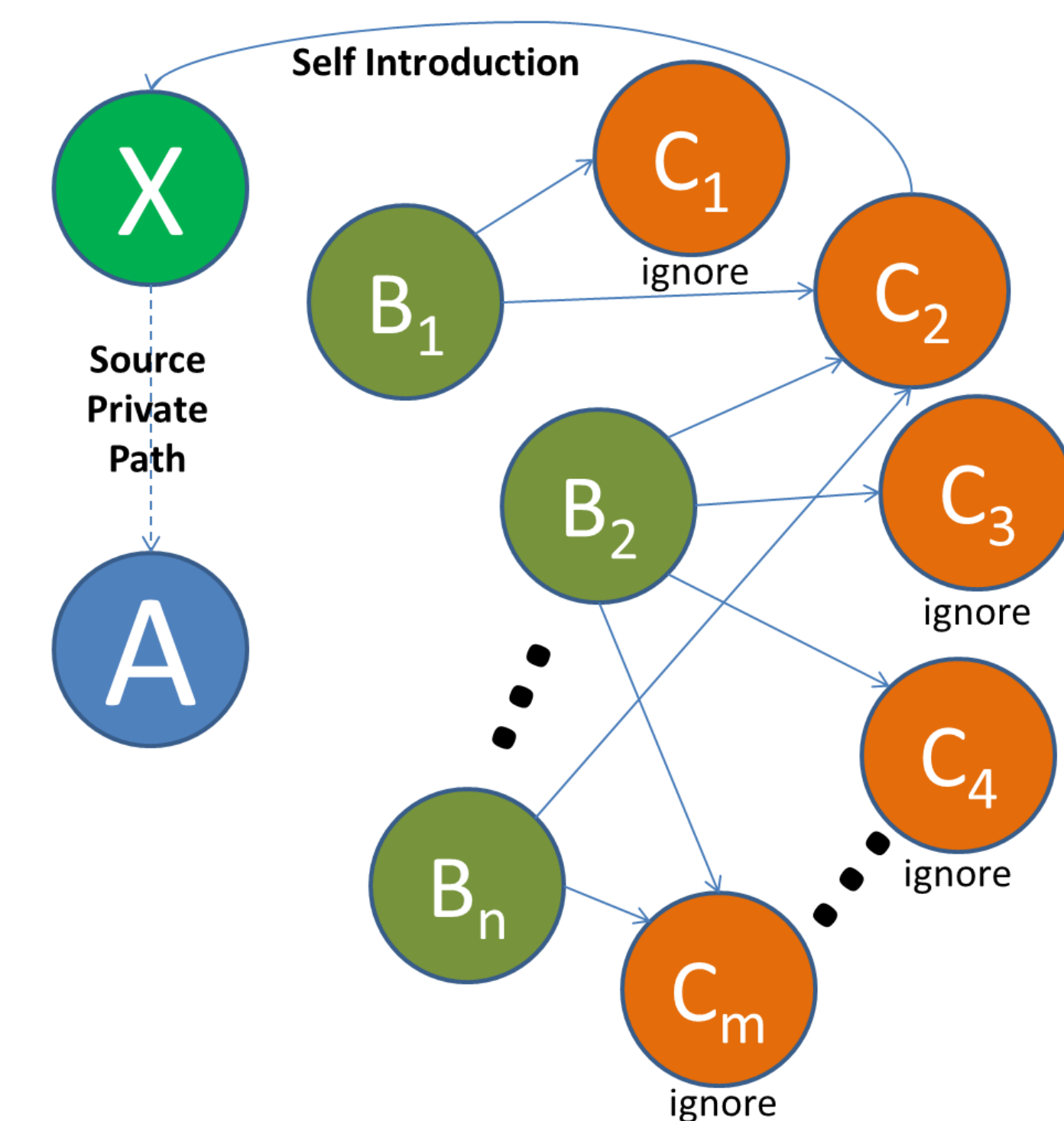
A trusted server T is used to public the encryption key and do the decryption. Only T knows the private key for decryption.

Each B_i encrypts his/her friend list as a vector and send it to A. A sums up all the list, permute the list and send it to T. T decrypts only the numbers with high frequency, and send it back to A. A undo the permutation and get the common neighbor recommendation.

A hash function is used to represent the friend list, resulting in approximate recommendation.



Protocol: $PPFR_h$



Protocol: $PPFR_{sp}$

The protocol $PPFR_{sp}$ utilizes **protecting the source privacy**[2] technique to let the candidates introduce themselves.

All A's friends B_i will send announcement to their friends C_j together with an randomly chosen message passing path.

C_j can choose different policy to decide whether send an self introduction along the given path.

A will get multiple messages from anonymous users. A can read the self introduction only when he/she get enough number of messages from the same C_j . This is ensured by use **secrete sharing technique**[3].

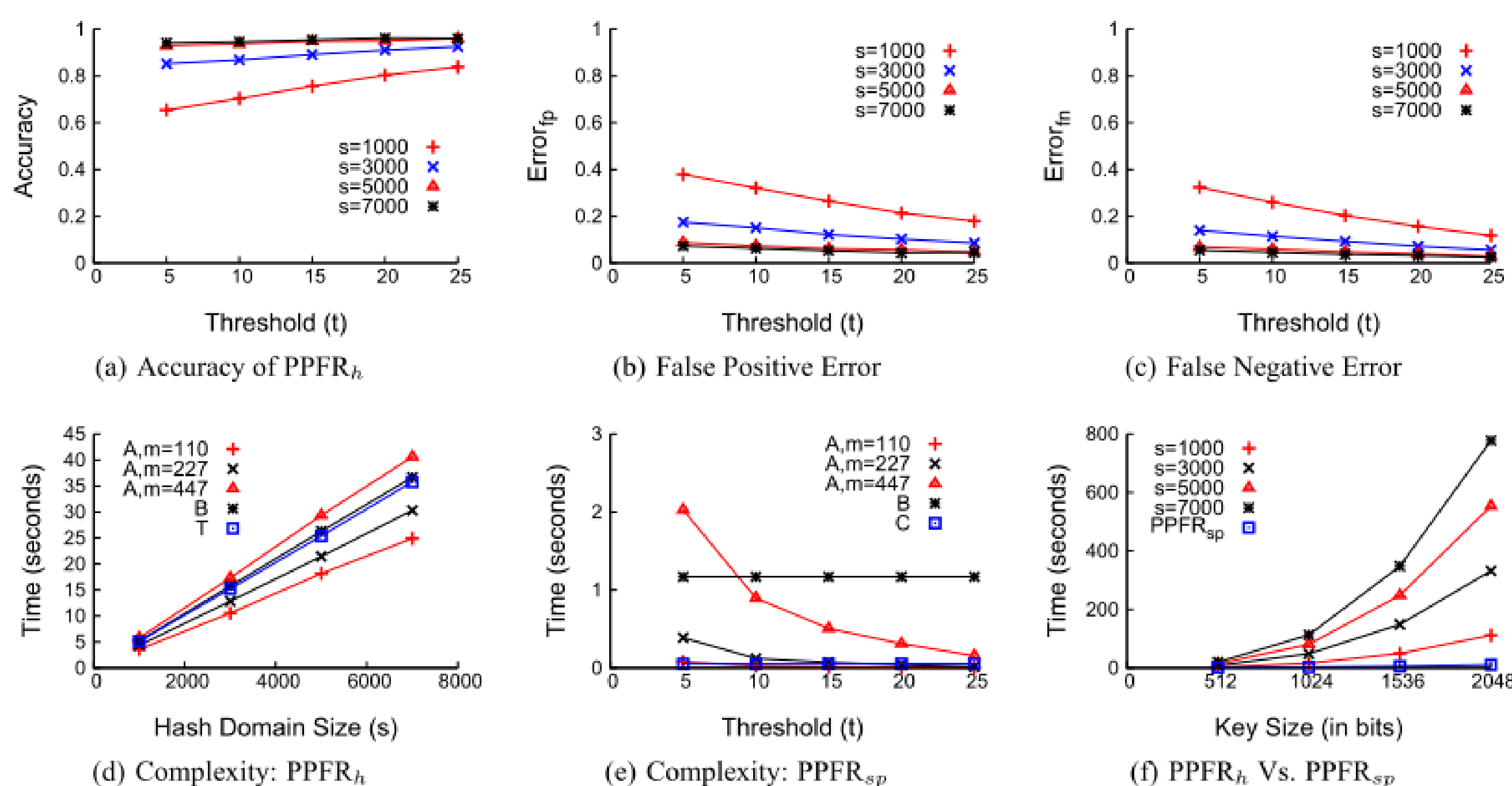
Public key system(RSA) is used to ensure source authentication.

Comparison

	$PPFR_h$	$PPFR_{sp}$
Computation Cost	$O(s \cdot Fr(A))$ s is the hash space size.	$O(Fr(B))$
Communication Cost	$O(K \cdot s \cdot Fr(A))$ K is the encryption key size.	$O(k \cdot l)$ k is the size of path. l is the number of (B_i, C_j) pairs.
Security (Semi-honest assumption)	More secure.	Less secure. A gets the exact number of common friends

In all, $PPFR_h$ is more secure under assumption and $PPFR_{sp}$ is more efficient and accurate. Our protocols act as a trade-off among security, efficiency, and accuracy thereby providing more flexibility to the users.

Empirical Results



[1] P. Paillier. Public-key cryptosystems based on composite degree residuosity classes. In *Proceedings of the 17th international conference on Theory and application of cryptographic techniques*, Springer-Verlag, 1999.

[2] W. Jiang, L. Si, and J. Li. Protecting source privacy in federated search. In *Proceedings of the 30th annual ACM SIGIR conference on Research and development in information retrieval*, 2007.

[3] A. Shamir. How to share a secret. *Communications of the ACM*, 22(11):612–613, 1979.