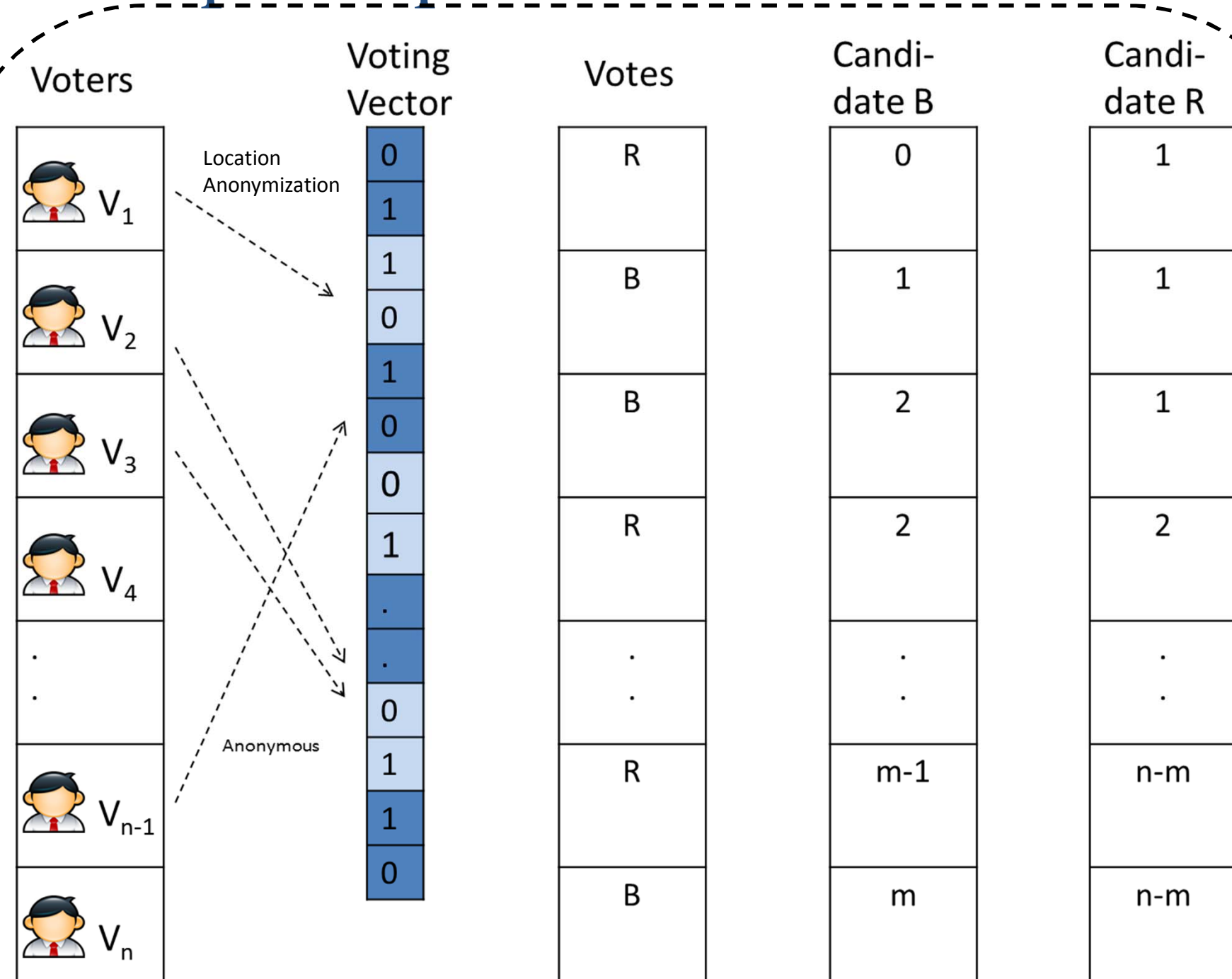


Mutual Restraining Voting Involving Multiple Conflicting Parties

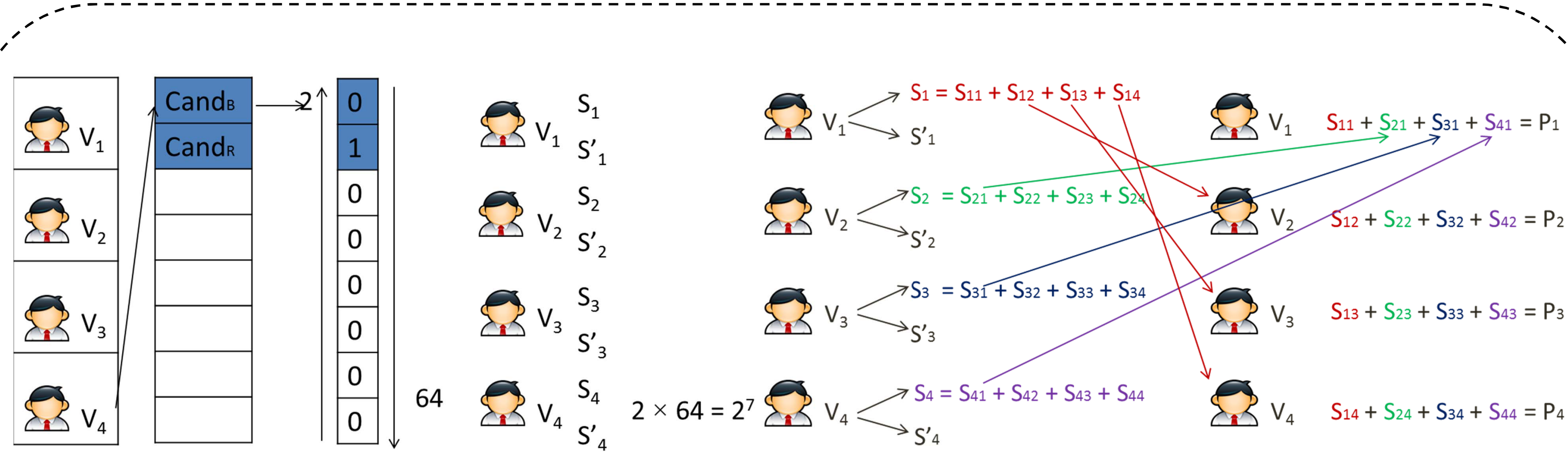
A. Problem Statement

- A gap between casting secret ballots and tallying & verifying individual votes.
- Due to disconnection between the vote-casting process and the vote-tallying process or opaque transition (e.g., due to encryption) from vote-casting to vote-tallying.
- A groundbreaking e-voting protocol that fills this gap and provides a fully transparent election.

B. Proposed protocol



TP1: Visual bulletin board: universally verifiable tallied **voting vector** and (incremental) tallies.

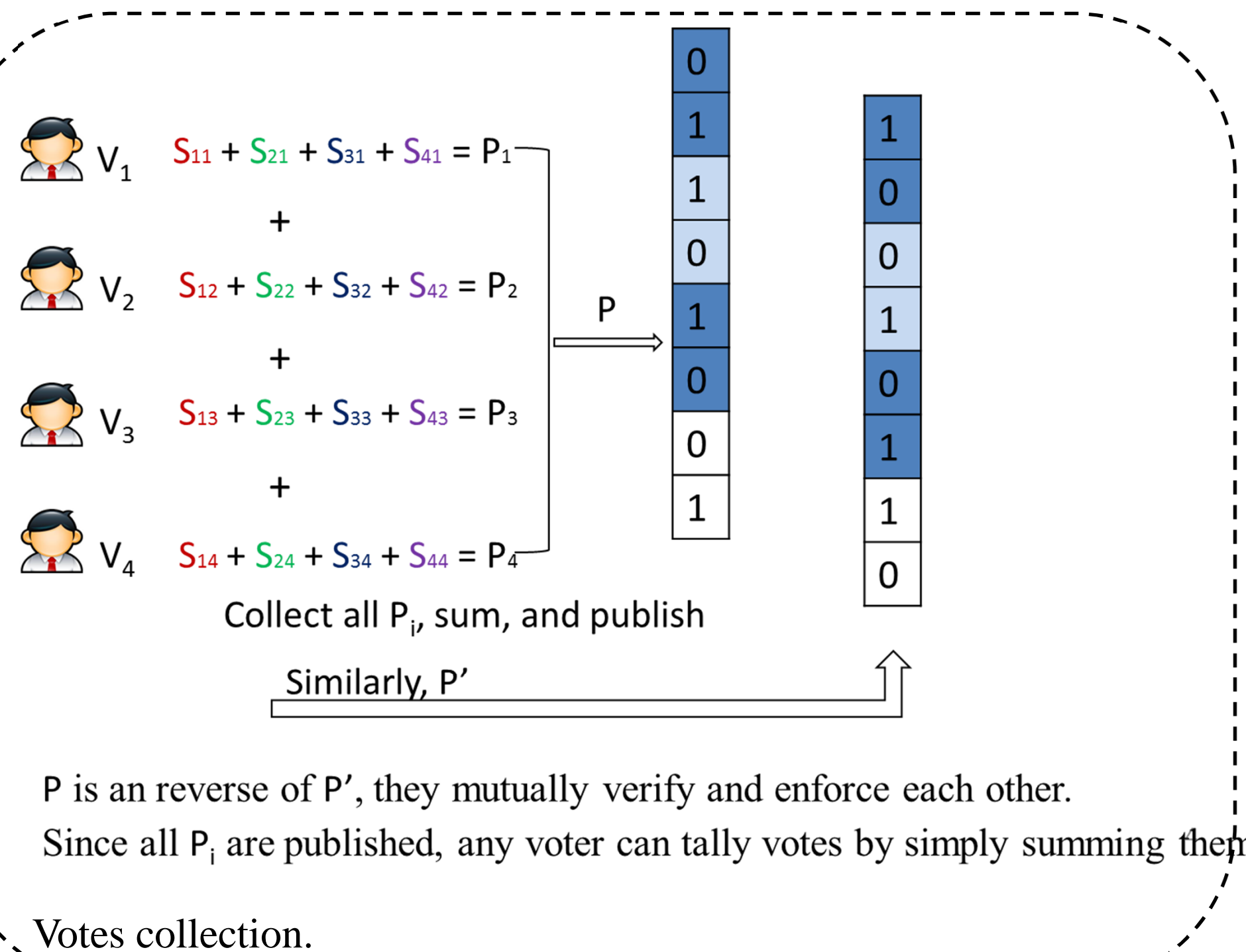


$$S_4 \times S'_4 = 2 \times 2^6 = 2^7$$

For any voter, $S_i \times S'_i = 2^{N-1}$ where N ($=4 \times 2=8$ here) is the length of the voting vector. Enforce one and just one vote.

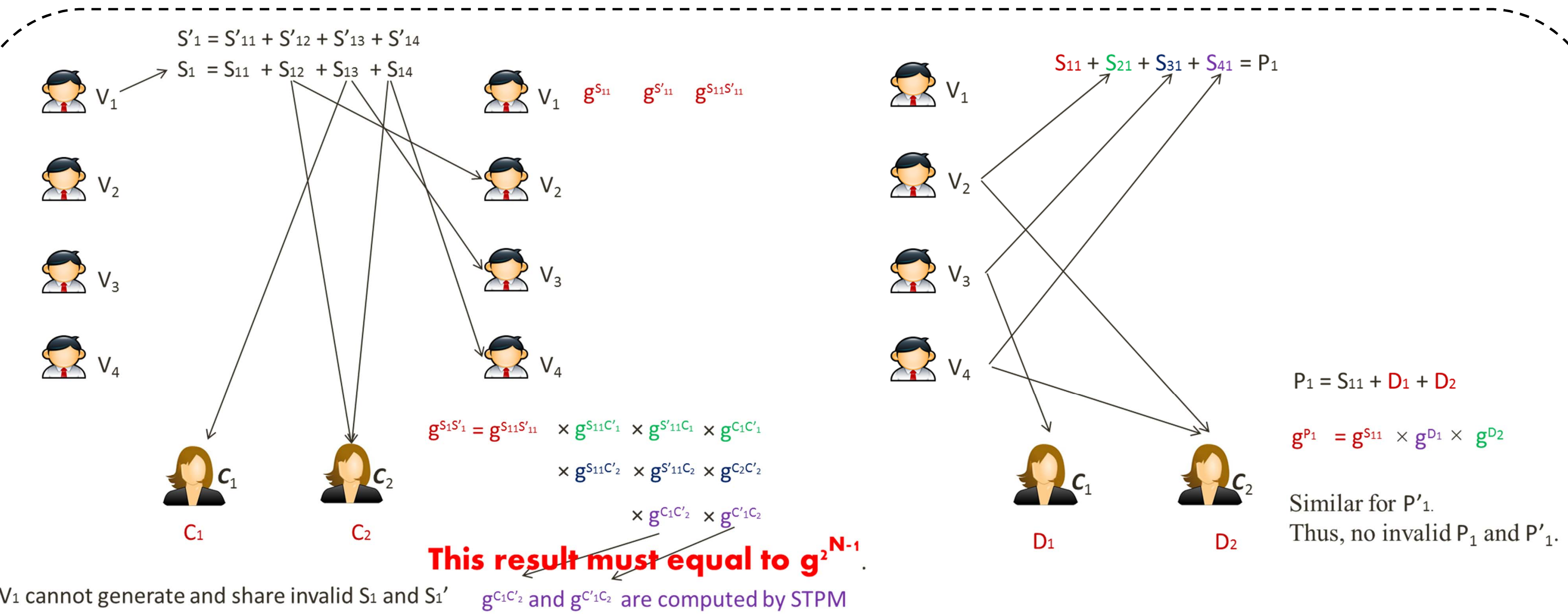
1. Split S_i into and distribute shares
 2. Sum up received shares (and its own) and publish the result
- Similarly, for S'_i .

TP2: forward and backward mutual lock voting by simplified (N,N) secret sharing



P is an reverse of P' , they mutually verify and enforce each other.
Since all P_i are published, any voter can tally votes by simply summing them.

Votes collection.



TP3: in-process verification and enforcement, using Secure Two Party Multiplication (STPM).

C. A Voting Example and Web Based Dynamic Bulletin Board

Voter	Secret Location	Vote	Shares			Secret ballot
			Self-computed	Server generated		
V ₁	2	B (32)	12 (=32-5-15)	<u>5</u>	8,7 (sum=15)	45 (=12+1+15+17)
V ₂	3	R (4)	13 (=4-1-(-10))	<u>1</u>	-3,-7 (sum=-10)	28 (=5+13+7+3)
V ₃	4	B (2)	-10 (=2-15-(-3))	15	<u>7,-10</u> (sum=-3)	30 (=8+(-3)+(-10)+35)
V ₄	1	R (64)	9 (=64-17-38)	17	<u>3,35</u> (sum=38)	-1 (=7+(-7)+(-10)+9)

A voting example involving 4 voters and 2 candidates (R and B): Notes: shares with underline are generated by Server 1, e.g., 5 of V_1 and 7 of V_3 , and **shares** in red are generated by Server 2, e.g., 8 of V_1 and 15 of V_3 .

What we get?

- Seamless, viewable, verifiable, and privacy-preserving transition from vote-casting to vote-tallying
- Individual voters can verify their own votes and are technically and visually assured that their votes are indeed counted in the final tally
- Public can verify the accuracy of the count, political parties will be able to catch fraudulent votes
- Secrecy of any voter's vote is remained
- Transparent e-voting protocol: enable open and fair elections with full voter assurance, even for the voters of minor or weak political parties.

Incremental aggregation			Incremental tallying			
Voter	Secret Ballot	Aggregation	V _A	Vote	R counts	B counts
V ₂	28	28	0	R	1	0
V ₁	45	73	1			
V ₄	-1	72	1	B	1	1
V ₃	30	102	0			
1.Incremental aggregation of the cast secret ballots			0	R	2	1
2.All partial aggregations 28, 73, and 72 has no information on votes			1			
3. Last aggregation 102 (=32+4+2+64) exposes all votes and it is the final tallied voting vector V _A			1	B	2	2
			0			

1. Incremental aggregation of the cast secret ballots
2. All partial aggregations 28, 73, and 72 has no information on votes
3. Last aggregation 102 ($=32+4+2+64$) exposes all votes and it is the final tallied voting vector V_A

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