

A Comprehensive Access Control System for Scientific Applications

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Problem Statement

- Web based scientific applications provide means to share scientific data beyond the local computing environment
- The organization and sharing of large and heterogeneous data pose challenges due to their sensitive nature

There is a need for a robust authorization mechanism to prevent unauthorized access to scientific data

- For this purpose, we present an access control system for scientific applications
- We formulate a methodology that incorporates principles from security management and software engineering

Authorization Requirements

- Implicit Authorization:** An explicitly specified authorization may imply authorizations i.e. authorizations can be automatically propagated
- Dataset Security:** A user having authorization to execute a tool should not have any authorization to directly modify the dataset accessed by the tool
- Sandbox Search:** A user is allowed to only execute a browsing query on the existence of data
- Temporal Constraints:** Permissions have a temporal dimension
- Conflict Resolution:** Identifying and resolving a conflict is essential in improving usability of any access control system.

Authorization Model

Basic Definition An authorization is defined as (s, o, p, s', c) where: $s \in S$, the set of subjects; $o \in O$, the set of objects; $p \in P$, the set of permission; $s' \in \text{owner}(o) \subseteq S$; $c \in C$, the set of class of objects. A function f is defined to determine if an authorization (s, o, p, s', c) is True or False;

$$f : S \times O \times P \times S \times C \rightarrow \{\text{True, False}\}$$

Definition 1. A positive authorization is a tuple (s, o, p, s', c) with $s \in S, o \in O, p \in P, s' \in S$ and $c \in C$. A negative authorization is a tuple $(s, o, \neg p, s', c)$ with $s \in S, o \in O, p \in P, s' \in S$ and $c \in C$.

Definition 2. An authorization base (AB) is a set of explicit authorizations (s, o, p, s', c) with $s \in S, o \in O, p \in P, s' \in S$ and $c \in C$ where p positive or negative; that is,

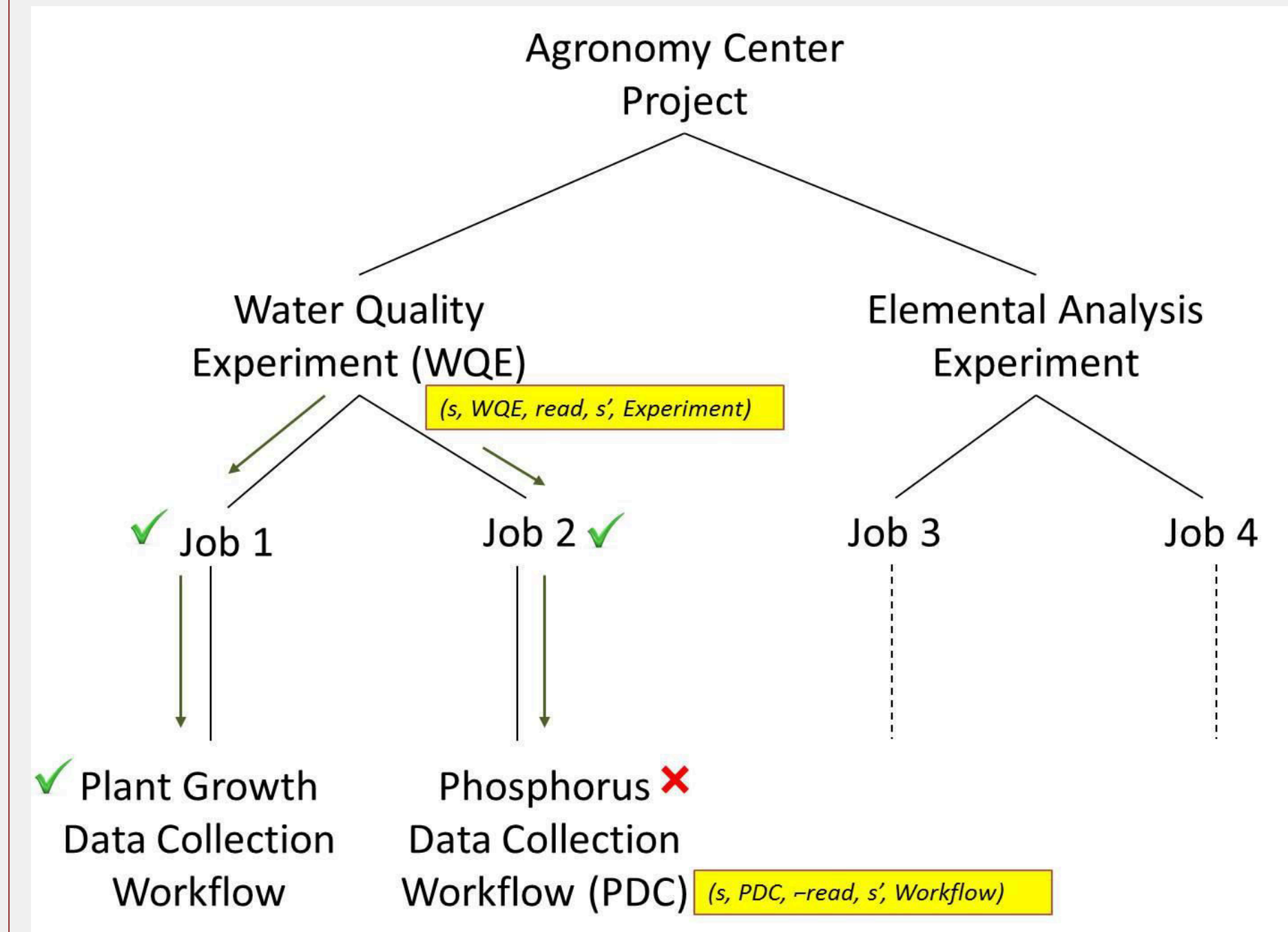
$$AB \subseteq S \times O \times P \times S \times C$$

Definition 3. Function $v(s, o, p, s', c)$ is defined as

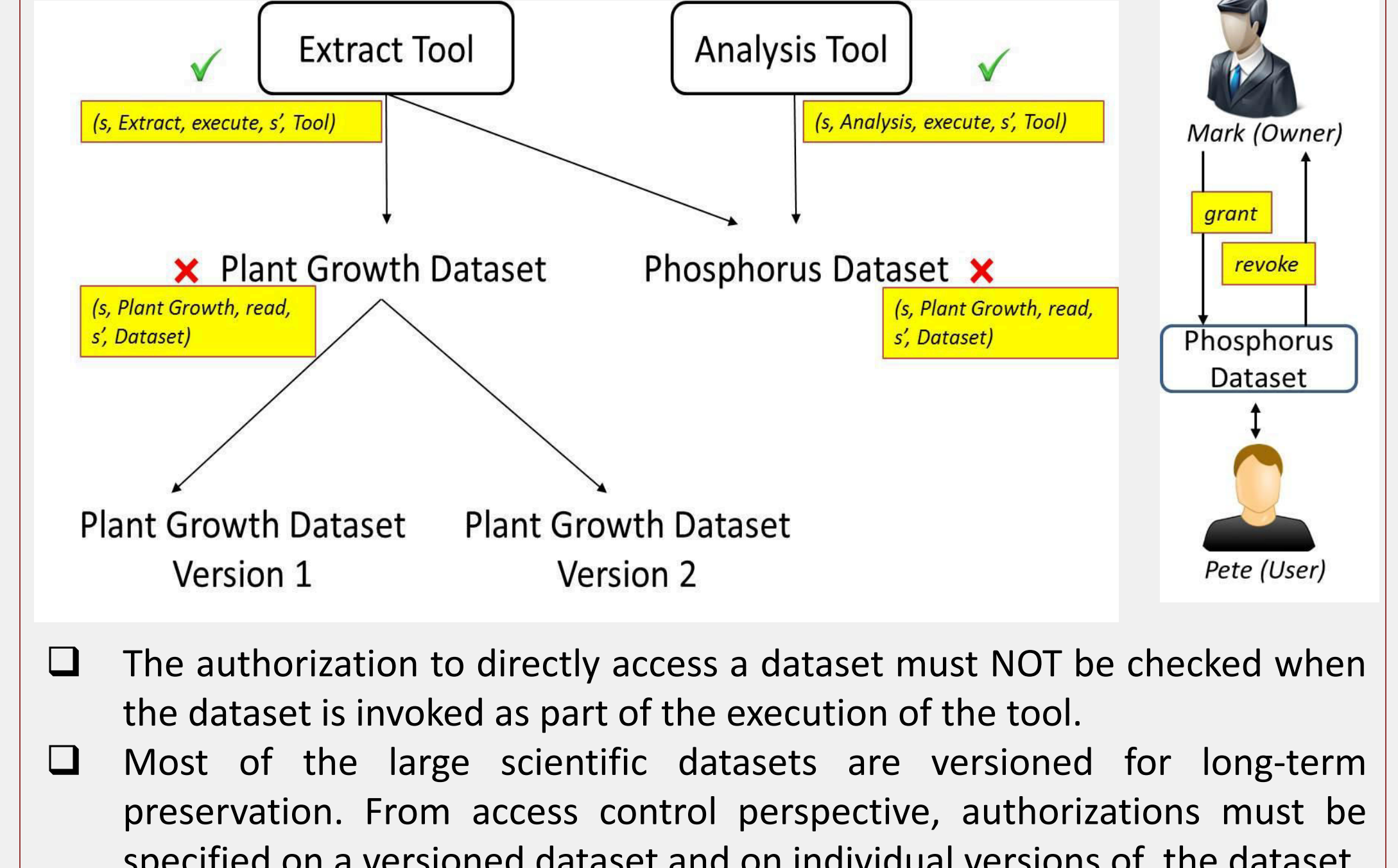
$$v : S \times O \times P \times S \times C \rightarrow \{\text{True, False}\}$$

If $(s, o, p, s', c) \in AB$, then $v(s, o, p, s', c) = \text{True}$; else, if there exists an $(s_1, o_1, p_1, s'_1, c_1) \in AB$ such that $(s_1, o_1, p_1, s'_1, c_1) \rightarrow (s, o, p, s', c)$, then $v(s, o, p, s', c) = \text{True}$; else, if there exists an $(s_1, o_1, \neg p_1, s'_1, c_1) \in AB$ such that $(s_1, o_1, \neg p_1, s'_1, c_1) \rightarrow (s, o, \neg p, s', c)$, then $v(s, o, p, s', c) = \text{False}$.

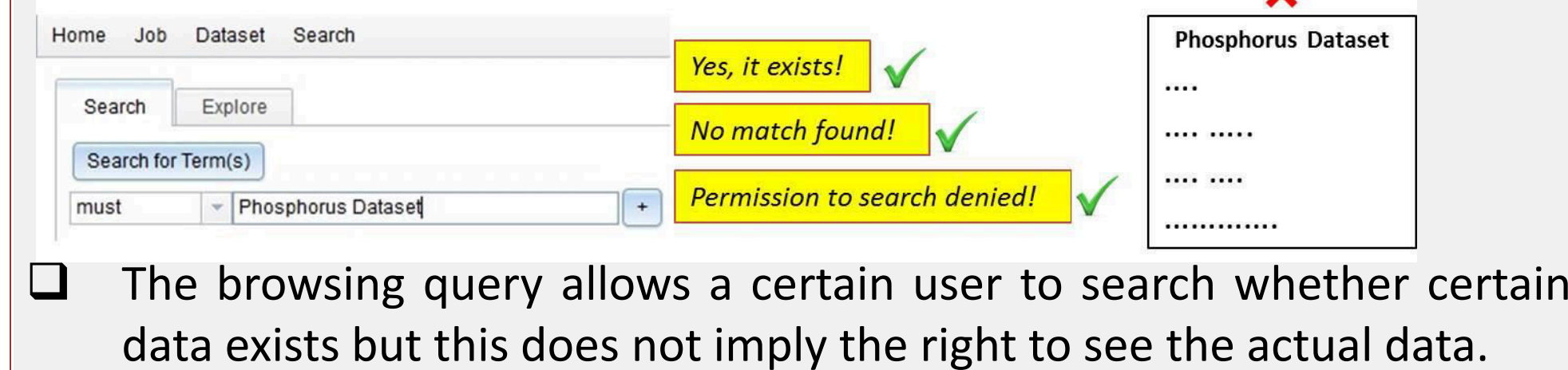
Implicit Authorization



Dataset Security



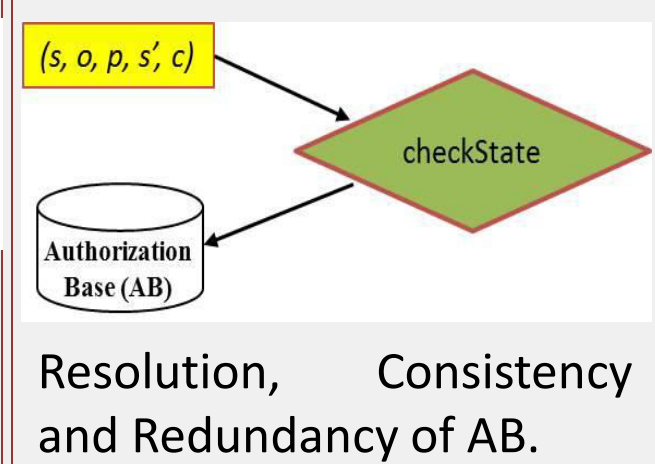
Sandbox Search



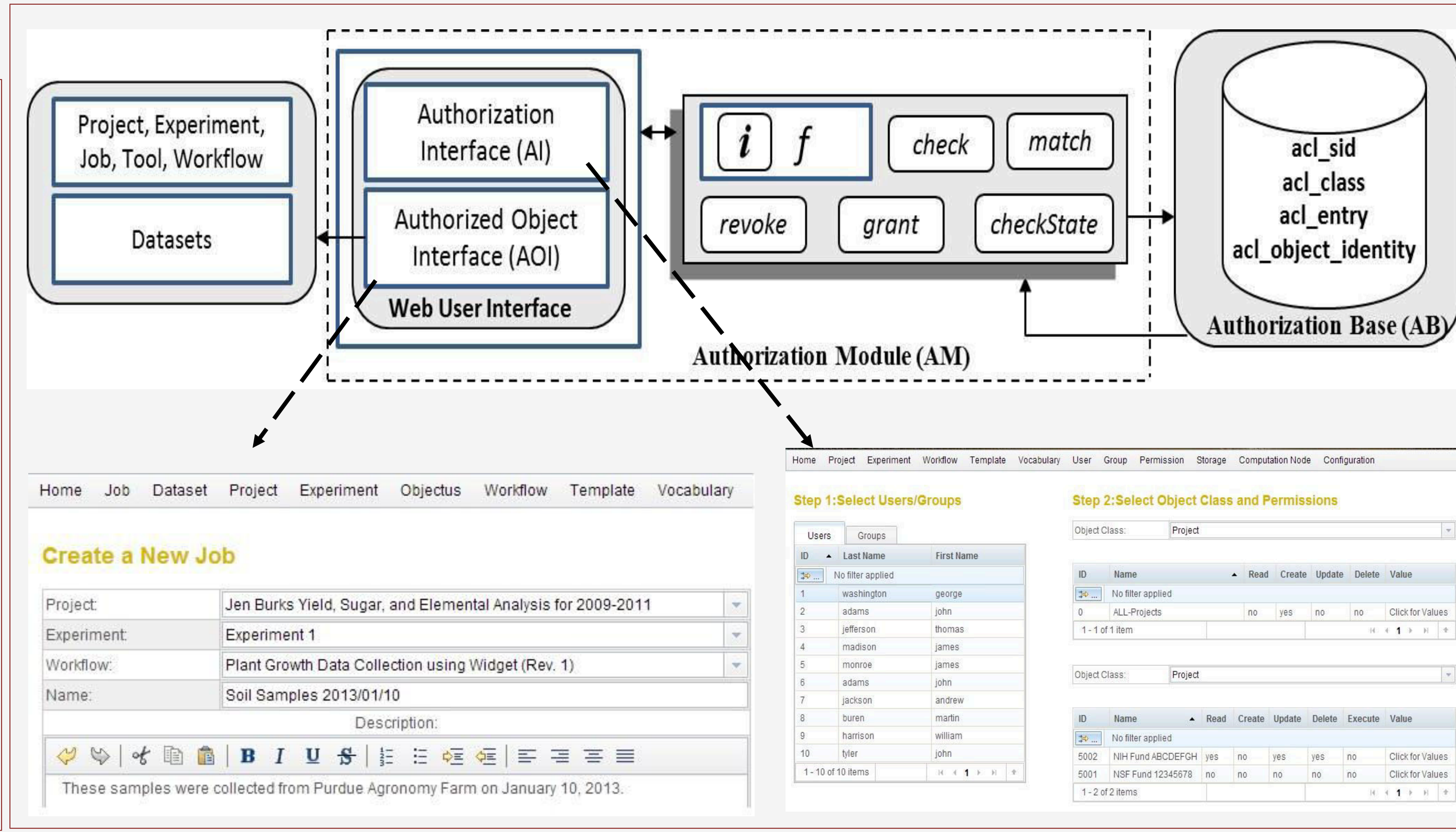
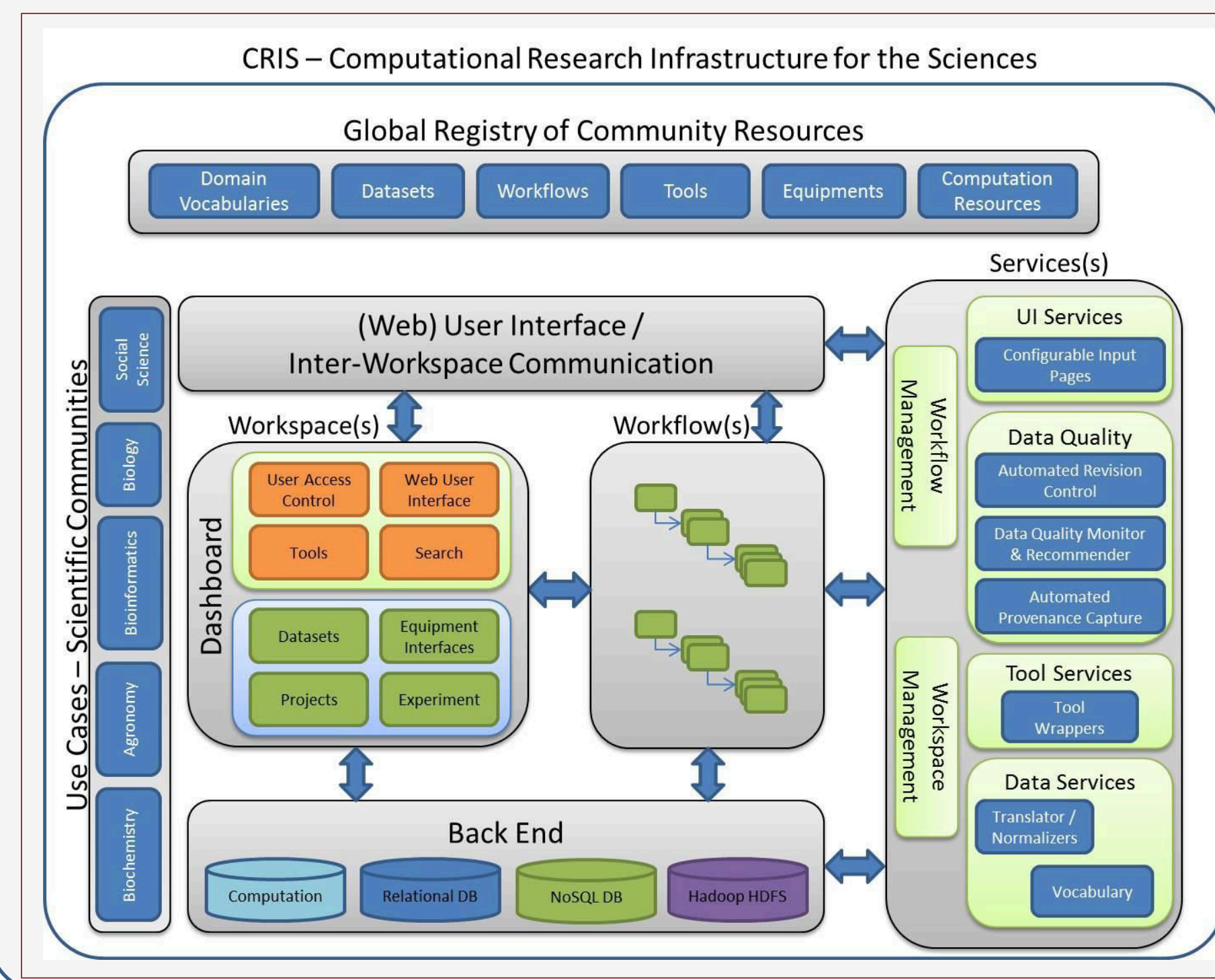
Temporal Constraint

Definition 4. A temporal authorization is a pair (period, auth), where period is a time interval $[t_a, t_b]$ with $t_a \in \mathbb{N}, t_b \in \mathbb{N} \cup \infty, t_a \leq t_b$, and $\text{auth} = (s, o, p, s', c)$.

Conflict Resolution



CRIS Access Control Architecture



- CRIS is a web based application with its primary tenets to provide an easy to use, scalable and collaborative scientific infrastructure for scientists.
- CRIS has been implemented using open source software and free Web APIs
- Since CRIS is Spring-based, we adopt the authorization modules provided by Spring Security as it is the de-facto standard for securing Spring-based applications.
- The AM in the architecture provides a CRIS user the ability to store/create authorizations through AI and consequently be allowed access to authorized objects through AOI.